

Minimal Overlapping & Exact Matches in Words

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A Quick Outline

- Basic Definitions
- Where This Problem Started
- What is the Problem?
- Results & Applications

Definitions

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Note: For this talk, we will concern ourselves with the alphabets $\mathbb{N} = \{1, 2, 3, \dots\}$ and $\mathbb{B} = \{0, 1\}$.

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Definition: $n_a(w)$ is the number of occurrences of a in w .

Motivation

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How many words in \mathbb{N}^ have weight n and contain exactly one occurrence of the word 13?*

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Example: ($n = 6$) There are 12 words - 15, 24, 114, 141, 33, 123,
132, 213, 231, 1113, 1131, 1311.

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Example: ($n = 6$) There are 12 words - 15, 24, 114, 141, 33, 123,
132, 213, 231, 1113, 1131, 1311.

Example: ($n = 7, 8$) There are 30 words of weight 7 and 68 of weight 8.

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Here we are talking about consecutive patterns (Sergi's talk, not Bridget's talk).

More Definitions

Let $w = w_1w_2 \cdots w_n \in \mathbb{N}^n$ and $u = u_1u_2 \cdots u_j \in \mathbb{N}^j$ with $j \leq n$.

Definition: w has an *exact u -match* starting at position i if

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Definition: w has an *endpoint embedding* of u starting at position i if $w_i \geq u_1$, $w_{i+j-1} \geq u_j$, and $w_{i+k} = u_{k+1}$ for $2 \leq k \leq j - 1$.

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Definition: $ex_u(w)$ is the number of exact u -matches in w , and $endp_u(w)$ is the number of endpoint embedding of u in w .

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So, $ex_u(w) = 2$ and $endp_u(w) = 3$.

More Definitions

Definition: $u \in \mathbb{N}^j$ has the *end point minimal overlapping property* if $i = n(j - 1) + 1$ is the smallest i such that there exists $w \in \mathbb{N}^i$ with $endp_u(w) = n$. We denote the set of such w 's by $\mathcal{EPM}\mathcal{P}_{u, n(j-1)+1}$, and we refer to the elements of this set as *end point embedding maximum packings of u* .

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Definition: $epmp_{u, n(j-1)+1}(z_0, z_1, \dots) = \sum_{w \in \mathcal{EPMP}_{u, n(j-1)+1}} z(w).$

A Computational Example

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 $w = w_1 32 w_4 32 \dots w_{3n-2} 32 w_{3n+1}$, where $w_1 \geq 5$, $w_{3n+1} \geq 4$,
and $w_{3k+1} \geq \max(4, 5)$ for $1 \leq k \leq n-1$. So,
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and $w_{3k+1} \geq \max(4, 5)$ for $1 \leq k \leq n-1$. So,

$$epmp_{5324,3n+1}(z_0, z_1, \dots) = \left(\sum_{i \geq 5} z_i \right) \left(\sum_{i \geq 4} z_i \right) \left(\sum_{i \geq \max(4,5)} z_i \right)^{n-1} (z_2 z_3)^n.$$

A Main Theorem

Theorem: Suppose $u \in \mathbb{N}^j$ has the end point minimal embedding property. Then

$$\sum_{n \geq 0} t^n \sum_{w \in \mathbb{N}^n} x^{\text{endp}_u(w)} z(w) =$$

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Proof: Brick tabloids, involution principle, and lots of cases.

More Definitions

Definition: $u \in \mathbb{N}^j$ has the *exact match minimal overlapping property* if $i = n(j - 1) + 1$ is the smallest i such that there exists $w \in \mathbb{N}^i$ with $ex_u(w) = n$. We denote the set of such w 's by $\mathcal{EM}\mathcal{P}_{u, n(j-1)+1}$, and we refer to the elements of this set as *exact match maximum packings of u* .

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Example: For any $k \in \mathbb{N}$, $u = 01^k0 \in \mathcal{EM}\mathcal{P}_{u, n(k+1)+1}$.

Another Main Result

Theorem: Suppose $u = u_1u_2 \cdots u_j \in \mathbb{N}^j$ with $j \geq 3$ with
 $\min(u_1, u_j) > \max(u_2, \dots, u_{j-1})$.

Another Main Result

Theorem: Suppose $u = u_1 u_2 \cdots u_j \in \mathbb{N}^j$ with $j \geq 3$ with $\min(u_1, u_j) > \max(u_2, \dots, u_{j-1})$. If $\Theta(u) = 0^{u_1} 10^{u_2} 1 \cdots 10^{u_j}$, then

$$\sum_{v \in \mathbb{B}^*} x^{\text{ex}_{\Theta(u)}(v)} t^{n_1(v)} z^{|v|} = \text{Some big, messy g.f.}$$

An Application

Example: Let $u = 312$, so that $\Theta(u) = 0^31010^2$, which does not have the minimal overlap property. Then we can use our previous theorem to find the following:

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Example: Let $\Delta = \{303, 313, 323\}$, so that
 $\Theta(\Delta) = \{0^3110^3, 0^31010^3, 0^310^210^3\}$. Then we can extend our
previous theorem to find the following:

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previous theorem to find the following:

$$\sum_{\substack{v \in \mathbb{B}^* \\ ex_{\Theta(\Delta)}(v)=0}} z^{|v|} = \frac{1 + z^5 + 2z^6 + 3z^7 + 2z^8 + z^9}{1 - 2z + z^5 - z^7 - 3z^8 - 2z^9 - z^{10}}.$$

Future Directions

Question: Is there anything special about binary words?

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Answer: Yes & No. (This is joint work with Tom Langley)

The End

Thanks for listening and thanks to all of the organizers.