The Fermat Factorization Method

R. C. Daileda



Trinity University

Number Theory

Introduction

Recall: The security of the RSA cryptosystem depends on the difficulty in factoring the encryption modulus n = pq.

Poor choices of p and q can lead to easily factored values of n, rendering the cryptosystem "cracked."

One such situation occurs when p and q are relatively close together. In this case one can apply the *Fermat Factorization Method* to find p and q.

Remark: The Fermat method can be applied to arbitrary odd n to try to find a divisor/complementary divisor pair that are relatively close together, if such a pair exists.

The Set-up

Suppose that n = ab with a > b odd. Notice that

$$n = ab$$

$$= \left(\frac{a+b}{2} + \frac{a-b}{2}\right) \left(\frac{a+b}{2} - \frac{a-b}{2}\right)$$

$$= \left(\frac{a+b}{2}\right)^2 - \left(\frac{a-b}{2}\right)^2.$$

If a and b are close together, then:

•
$$\frac{a-b}{2}$$
 is relatively small; specifically we assume $\frac{a-b}{2\sqrt{2b}} < \epsilon$.
• $\frac{a+b}{2}$ is not much larger than \sqrt{n} .

To quantify this final statement note that

$$n = \left(\frac{a+b}{2}\right)^2 - \left(\frac{a-b}{2}\right)^2 \Rightarrow \left(\frac{a+b}{2}\right)^2 - n = \left(\frac{a-b}{2}\right)^2$$

and hence

$$\left(\frac{a+b}{2}-\sqrt{n}\right)\left(\frac{a+b}{2}+\sqrt{n}\right)=\left(\frac{a-b}{2}\right)^2.$$

Since

$$\left(\frac{a+b}{2}+\sqrt{n}\right)>2b$$
 and $\left(\frac{a-b}{2}\right)^2>0$

we obtain

$$0 < \frac{a+b}{2} - \sqrt{n} < \left(\frac{a-b}{2\sqrt{2b}}\right)^2 < \epsilon^2.$$

The Algorithm: Fermat Factorization

Moral: If *n* is the product of two distinct odd numbers that are close together, then

$$n=t^2-s^2$$

where t is slightly larger than \sqrt{n} and s is relatively small.

How can we use this to factor n? Set $t_0 = \lceil \sqrt{n} \rceil$ and successively compute

$$\sqrt{t_0^2-n}, \sqrt{(t_0+1)^2-n}, \sqrt{(t_0+2)^2-n}, \sqrt{(t_0+3)^2-n}, \dots$$

until one obtains

$$\sqrt{(t_0+k)^2-n}=s\in\mathbb{N}.$$

If we set $t = t_0 + k$, then $n = t^2 - s^2 = (t + s)(t - s)$.

Remarks:

- Because of our assumption on (a b)/2, this process is guaranteed to stop after roughly ϵ^2 steps.
- The factors (t + s) and (t s) are nontrivial because $t \sim \sqrt{n}$ while $s \sim 0$.

Example

Apply the Fermat Factorization Method to factor

n = 2251644881930449333.

We have

$$t_0 = \lceil \sqrt{n} \rceil = 1500548194.$$

Moreover, we find that

$$\begin{split} \sqrt{t_0^2 - n} &= \sqrt{586212303} = 24211.821\dots, \\ \sqrt{(t_0 + 1)^2 - n} &= 2\sqrt{896827173} = 59894.145\dots, \\ \sqrt{(t_0 + 2)^2 - n} &= \sqrt{6588405083} = 81168.990\dots, \\ \sqrt{(t_0 + 3)^2 - n} &= 97926 \end{split}$$

so that $t = t_0 + 3 = 1500548197$ and s = 97926. Hence n = pq with

$$p = t + s = 1500646123,$$

 $q = t - s = 1500450271,$

both of which turn out to be prime. No

te
$$\left(\frac{p-q}{2\sqrt{2q}}\right)^2 = 3.1955\ldots$$

Example

The integer

$n = 89564941429129460494158838187124492462610412156204 \\ 2227318384494381723497514540860474803494041479529$

is the product of two primes. Use Fermat Factorization to find them.

We have

 $t_0 = \lceil \sqrt{n} \rceil$

= 29927402397991286489627871143011285937749436382209

and with the aid of a computer we find that

$$\sqrt{(t_0+18)^2-n}=33408832099552561140000000.$$

Hence

- s = 33408832099552561140000000,
- t=29927402397991286489627871143011285937749436382227,

so that the prime factorization of n is pq where

$$\begin{split} p &= 29927402397991286489627837734179186385188296382227, \\ q &= 29927402397991286489627904551843385490310576382227. \end{split}$$

Note that
$$\left(\frac{p-q}{2\sqrt{2p}}\right)^2 = 18.6476\ldots$$

Remark: Factoring n in this manner only took a matter of minutes using Maple. However, after over 8 hours neither Maple nor PARI could successfully factor n naïvely.