

# Associativity of the Symmetric Difference

R. C. Daileda

Given sets  $A$  and  $B$ , their *symmetric difference* is

$$A\Delta B = (A \setminus B) \cup (B \setminus A) \quad (1)$$

$$= (A \cup B) \setminus (A \cap B). \quad (2)$$

Because (1) (and (2)) is symmetric in  $A$  and  $B$ , we immediately find that  $\Delta$  is commutative. That is,  $A\Delta B = B\Delta A$ . The purpose of this note is to prove the following less obvious property of the  $\Delta$  operation.

**Proposition 1.** *The symmetric difference is associative. That is, given sets  $A$ ,  $B$  and  $C$ , one has*

$$(A\Delta B)\Delta C = A\Delta(B\Delta C).$$

This proposition is an almost immediate consequence of the characterization of  $(A\Delta B)\Delta C$  given below.

**Lemma 1.** *Let  $A$ ,  $B$  and  $C$  be sets. Then*

$$(A\Delta B)\Delta C = \left( (A \cup B \cup C) \setminus [(A \cap B) \cup (A \cap C) \cup (B \cap C)] \right) \cup (A \cap B \cap C). \quad (3)$$

To see how the proposition follows from the lemma, note that the right hand side of (3) is invariant under permutation of  $A$ ,  $B$  and  $C$ . Thus

$$(A\Delta B)\Delta C = (B\Delta C)\Delta A = A\Delta(B\Delta C),$$

where we have used the commutativity of  $\Delta$  to obtain the final equality. So all we need to do now is prove the lemma.

*Proof of Lemma 1.* By (1), (2) and the distributive laws for  $\wedge$  and  $\vee$  we have<sup>1</sup>

$$\begin{aligned} x \in (A\Delta B)\Delta C &\cong (x \in A\Delta B \wedge x \notin C) \vee (x \notin A\Delta B \wedge x \in C) \\ &\cong \left( [(x \in A \wedge x \notin B) \vee (x \notin A \wedge x \in B)] \wedge x \notin C \right) \vee \left( [x \notin A \cup B \vee x \in A \cap B] \wedge x \in C \right) \\ &\cong [x \in A \wedge x \notin B \wedge x \notin C] \vee [x \notin A \wedge x \in B \wedge x \notin C] \vee [x \notin A \cup B \wedge x \in C] \vee [x \in A \cap B \wedge x \in C] \\ &\cong [x \in A \wedge x \notin B \wedge x \notin C] \vee [x \notin A \wedge x \in B \wedge x \notin C] \vee [x \notin A \wedge x \notin B \wedge x \in C] \vee [x \in A \cap B \cap C] \\ &\cong [x \in A \wedge x \notin B \cup C] \vee [x \in B \wedge x \notin A \cup C] \vee [x \in C \wedge x \notin A \cup B] \vee [x \in A \cap B \cap C]. \end{aligned}$$

For arbitrary statements  $P$  and  $Q$ , notice that

$$(P \vee Q) \wedge \neg Q \cong (P \wedge \neg Q) \vee (Q \wedge \neg Q) \cong P \wedge \neg Q,$$

---

<sup>1</sup>Because it is symmetric in  $A$ ,  $B$  and  $C$ , the final statement in this chain of equivalences is sufficient to prove Proposition 1. It needs to be manipulated further, however, to prove Lemma 1.

